

#### **31st ACM/SIGAPP Symposium on Applied Computing**

# SMT-Based Context-Bounded Model Checking for CUDA Programs

Phillipe Pereira, Higo Albuquerque, Hendrio Marques, Isabela Silva, Vanessa Santos, Celso Barbosa, Ricardo Ferreira, and **Lucas Cordeiro** 

Developed by NVIDIA to configure GPUs Seven Control Contro Control Control Control Control Control Control Control Contro





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- initially used in graphical processing in games applications

> specially those that require **high** computational power





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- Currently used in:
  - ➢ biomedicine
  - $\triangleright$  air traffic control
  - $\succ$  weather simulation





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- Currently used in:
  - ➢ biomedicine
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- We need to ensure code correctness in safety-critical GPU applications





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- Apply context-bounded model checking based on the Satisfiability Modulo Theories (SMT)
  - Monotonic Partial Order Reduction (MPOR) (CAV'09)
- Compare ESBMC-GPU experimental results with other state-of-art software verifiers for CUDA

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  - **Abstractly** represent the associated CUDA libraries
    - checks pre- and post-conditions
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  - Other extensions to ESBMC based on operational models
    - ESBMC++ (ECBS'13) and ESBMC<sup>QtOM</sup> (SPIN'16)
- CUDA is a proprietary platform
  - CUDA Programming Guide and IDE Nsight









```
#include <cuda.h>
#include <stdio.h>
#define N 2
  _global___ void definitions(int* A){
        atomicAdd(A,10);
int main (){
        int a = 5;
        int *dev_a;
        cudaMalloc ((void**) &dev_a, sizeof(int));
        cudaMemcpy(dev_a, &a,
sizeof(int),cudaMemcpyHostToDevice);
        ESBMC_verify_kernel(definitions,1,N,dev_a);
        cudaMemcpy(&a,dev_a,sizeof(int),cudaMemcpyD
eviceToHost);
        assert(a==25);
        cudaFree(dev_a);
        return 0;
}
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# cudaMalloc
cudaError_t cudaMalloc(void ** devPtr, size_t size) {
       cudaError_t tmp;
          ESBMC_assert(size > 0, "Size to be allocated must be greater than zero");
        *devPtr = malloc(size);
        if (*devPtr == NULL) {
                 tmp = CUDA ERROR OUT OF MEMORY;
                 exit(1);
         } else {
                 tmp = CUDA SUCCESS;
          _ESBMC_assert(tmp == CUDA_SUCCESS, "Memory was not allocated");
        lastError = tmp;
        return lastError;
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#### simulate behavior

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  - Using the Pthread/POSIX library

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#### **CUDA** program

_global_ void kernel(){ A[tidx.x]=tidx.x; }	
int main(){	
int *a; int *dev_a;	
cudaMalloc(&dev_a,a,size);	
cudaMemcpy(dev_a,a,htd);	
 ESBMC_verify_kernel( kernel,M,N,dev_a);	
cudaMemcpy(a,dev_a,dth);	
cudaFree(dev_a);	
rree(a);	
}	

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CUDA program	COM	
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free(a); }		

- Verification model adopts the CPU parallel processing

   Using the Pthread/POSIX library
  - COM **CUDA** program ESBMC\_verify\_kernel global void kernel(){ (kernel,M,N,dev a) Function conversion A[tidx.x]=tidx.x; cudaMalloc(&dev a,size) kernel<<<M.N>>> int main(){ int \*a; int \*dev a; gridDim = dim3(M);cudaMalloc(&dev a.a.size): assert(size>0); blockDim = dim3(N);\*dev a=malloc(size); cudaMemcpy(dev a,a,htd); if(\*dev a==NULL) dim3 conversion ESBMC \_verify\_kernel( kernel, M, N, dev a); exit(1);struct dim3: cudaMemcpy(a,dev a,dth); gridDim.x=M; blockDim.x=N; cudaFree(dev a); gridDim.y=1; blockDim.y=1; free(a); gridDim.z=1; blockDim.z=1;

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## Monotonic Partial Order Reduction (MPOR)

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- MPOR classifies thread transitions in a multithreaded program
  - Each transition may be dependent or independent
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- First application of the technique to verify CUDAbased programs
  - Reduction in time and verification effort
  - Elimination of threads interleavings that access different array positions

# **MPOR Applied to CUDA-based Programs**

- MPOR algorithm in the ESBMC-GPU
- **1.** function MPOR (v,  $\pi$ )
- 2. Check whether  $s_i$  exists in  $\pi$ ; otherwise, go to step 4
- 3. Check whether  $A_i$  produces a new state in  $\pi$ ; otherwise, go to step 5
- 4. Analyze whether  $\gamma(s_{i-1}, s_i)$  is independent on  $\pi$ ; otherwise, go to step 6
- 5. Return "independent" on  $\pi$  and terminates
- 6. Return "dependent" on  $\pi$  and terminates
- 7. end function
- MPOR algorithm in the ESBMC-GPU
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- $\mathbf{v} = (A_i, C_i, s_i)$   $A_i$ : active thread  $C_i$ : context switch  $s_i$ : current state
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kernel1(int *a){
a[threadIdx.x] = threadIdx.x;
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v_0: t_0, 0, a[0] = 0, a[1] = 0
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#### 7. end function

```
kernel (int *a)
if(a[1]==1)
a[threadIdx.x+2] = threadIdx.x;
else
a[threadIdx.x] = threadIdx.x;
```

 $v_0: t_0, 0, a[0] = 0, a[1] = 0$ 

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This technique is also used by other GPU kernel verification tools (*e.g.*, GPUVerify and PUG)

#### Fermi - Stream Multiprocessor

Instruction Cache						
Scheduler		Scheduler				
Dispatch		Dispatch				
Register File						
Core	Core	Core	Core			
Core	Core	Core	Core			
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*Warp* (32 threads)

Warp



#### Fermi - Stream Multiprocessor








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  - CUDA intrinsic variables (*e.g., uint4*)

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<b>Result\Tool</b>	ESBMC- GPU	GKLEE	GPUVERIFY	PUG	CIVL
True Correct	60	53	58	39	23
False Correct	67	56	56 30		24
True Incorrect	1	14	9	7	0
False Incorrect	3	7	8	11	3
Not supported	23	24	49	82	104
Time(s)	811	128	147	12	158

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contain errors	contain errors				0
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Total number of which the error i	benchmar	ks in gram	9	7	0
was found and a was reported	n error pa	th	8	11	3
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but the verifier did not find it			49	82	104
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