

ESBMC v7.7 : Automating Branch Coverage Analysis Using CFG-Based Instrumentation and SMT Solving

Author: Chenfeng Wei Chenfeng.wei@machester.ac.uk Co-Authors: Tong Wu, Rafael Sa Menezes, Fedor Shmarov, Fatimah Aljaafari, Sangharatna Godboley, Kaled Alshmrany, Rosiane de Freitas, Lucas C. Cordeiro





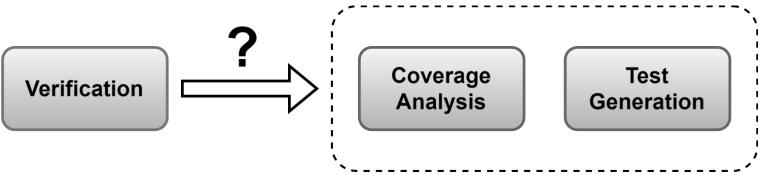




ESBMC, a bounded model checking (BMC) verifier based on SMT solving, has proven its effectiveness in bug detection in recent competitions.

However, it has never participated in the evaluation of the Cover-Branches category due to the lack of:

- Branch coverage analysis
- Automated test suite generation

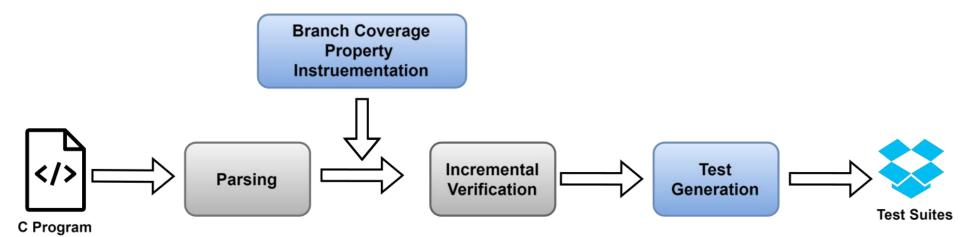


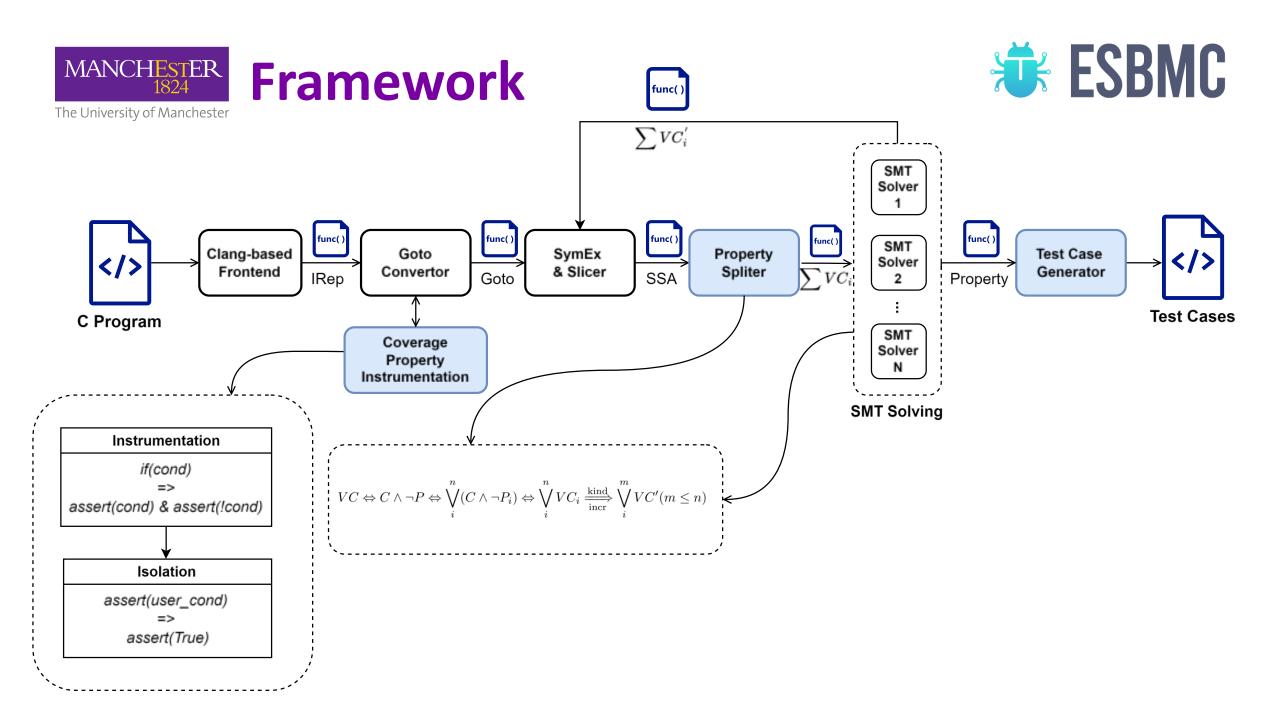


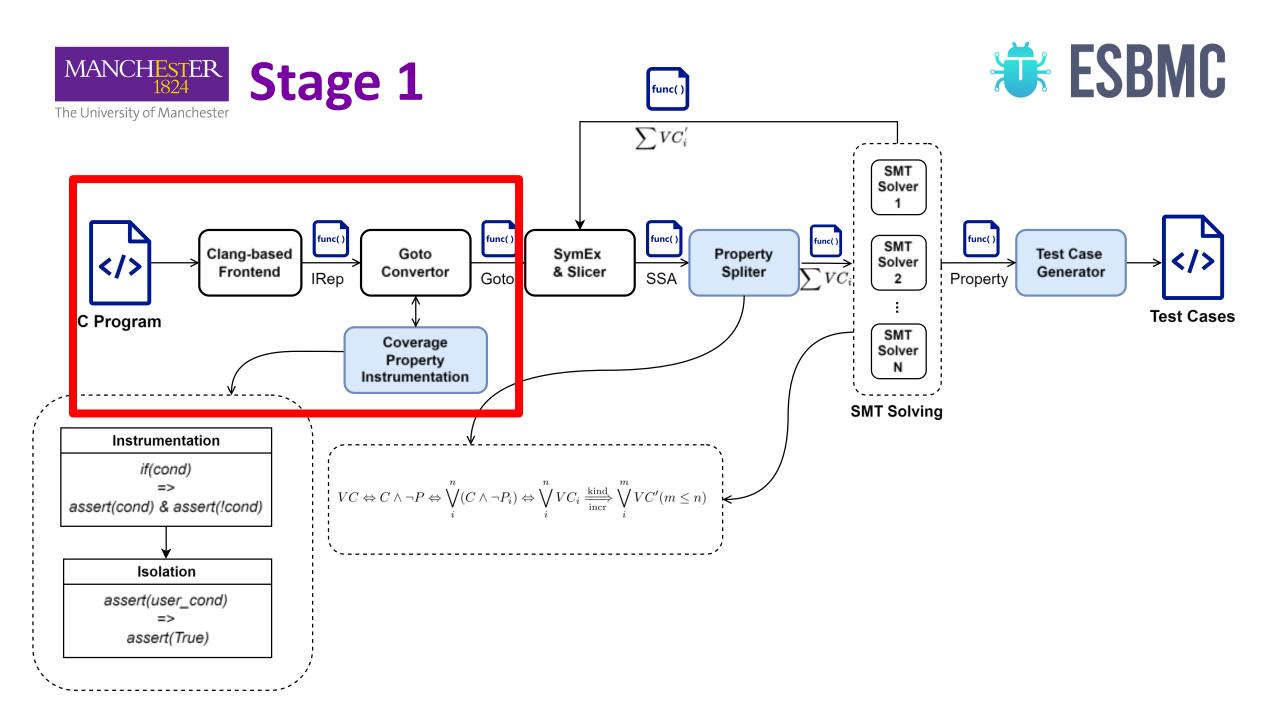


To bridge this gap, we present **ESBMC v7.7** with the following key contributions:

- Branch Coverage Instrumentation
 - Instrumentation & Isolation
- Incremental Multiple Property Verification
 - Property Splitting & Incremental Reasoning
 - Re-verification & Termination
 - Test Generation



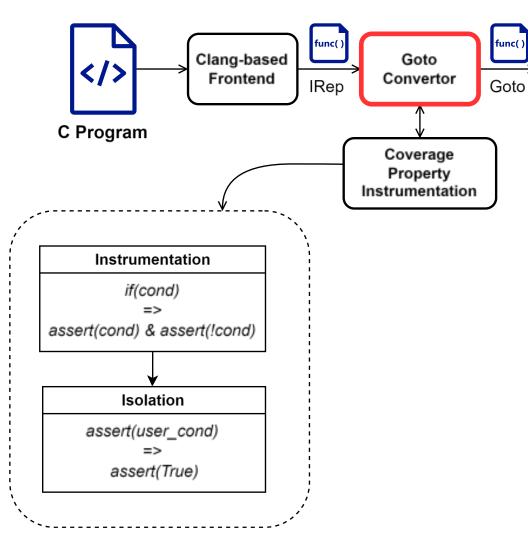




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func()





We apply our algorithm upon **GOTO program**, i.e., Control Flow Graph (CFG).

Simplification: In the Goto program, control-flow constructs such as if-else statements, while loops, for loops, and switch-case statements are normalized into if-goto structures.

```
int main()
                                                                        esbmc exp.c --goto2c &> exp2.c
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                                                 int x = VERIFIER nondet int();
                                               ESBMC goto label 1:;
int main()
                                                 Bool return value VERIFIER nondet bool 1;
                                                 return value VERIFIER nondet bool 1 = VERIFIER nondet bool();
  int x = 0;
                                                 if (!return value VERIFIER nondet bool 1) // while (nondet bool())
  while (__VERIFIER_nondet_bool())
                                                     goto ESBMC goto label 4;
                                                                                            // if (!x)
                                                 if (!(!( Bool)x))
                                                     goto ESBMC goto label 2;
    if (!x)
                                                 assert(x == 0);
                                                 assert(!overflow("+", x, 1));
                                                                                            // overflow-check
      assert(x == 0);
                                                 x = x + 1;
      ++X;
                                                 goto ESBMC goto label 3;
                                               ESBMC goto label 2:;
    else if (x == 1)
                                                 if (!(x == 1))
                                                                                            // else if (x == 1)
                                                     goto ESBMC goto label 3;
                                                 assert(x > 0);
      assert(x > 0);
                                                 x = 2;
      x = 2;
                                               ESBMC goto label 3:;
                                                 goto ESBMC goto label 1;
                                               ESBMC_goto_label_4:;
                                                 VERIFIER_assert((int)(x == 2));
  VERIFIER assert(x == 2);
```

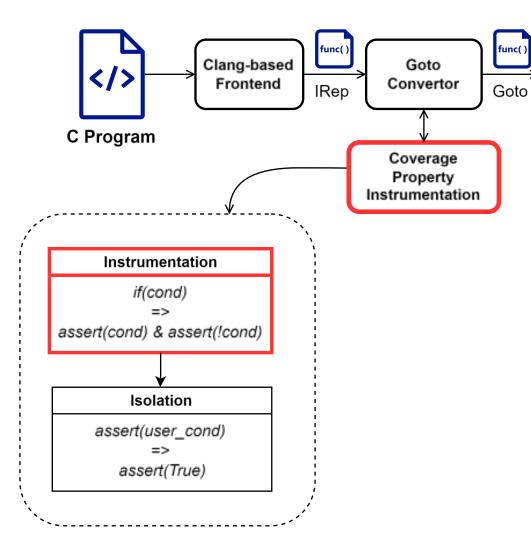
Instrumentation



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Instrumentation: to entry a branch guarded by condition like if(cond), there must exist an assignment that satisfies cond. This is equivalent to checking if a counterexample satisfies assert(!cond)

- True Branch: the body executes
- False Branch: the body is skipped

if(cond) => True branch: assert(!cond)
 => False branch: assert(!(!cond))





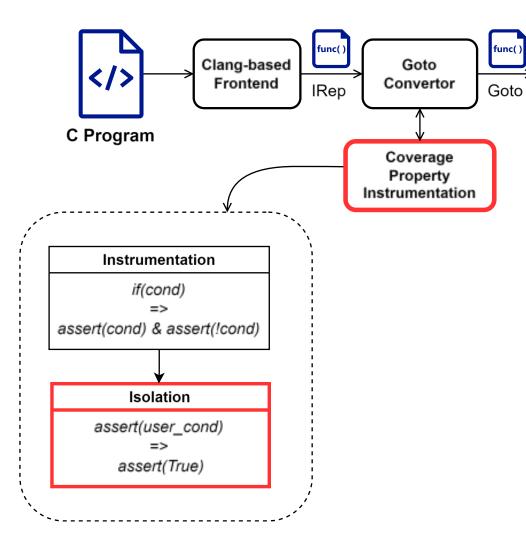
```
Bool return value VERIFIER nondet bool 1;
return value VERIFIER nondet bool 1 = VERIFIER nondet bool();
assert(!return_value____VERIFIER_nondet_bool_1); // !return_value$___VERIFIER_nondet_bool$1
assert(!(!return value VERIFIER nondet bool 1)); // !(!return value$ VERIFIER nondet bool$1)
if (!return value VERIFIER nondet bool 1)
   goto ESBMC goto label 4;
```

```
assert(!(!( Bool)x)); // !(!( Bool)x)
assert(!(!(!( Bool)x))); // !(!(!( Bool)x))
if (!(!( Bool)x))
   goto ESBMC goto label 2;
```

```
assert(!(x == 1)); // !(x == 1)
assert(!(!(x == 1))); // !(!(x == 1))
if (!(x == 1))
   goto __ESBMC_goto_label_3;
```







Isolation: Potential interferences are excluded to isolate the analysis of instrumented coverage properties from others.

User-defined properties, i.e. assertions, are converted into tautologies

assert(a[i]>0) => assert(True)

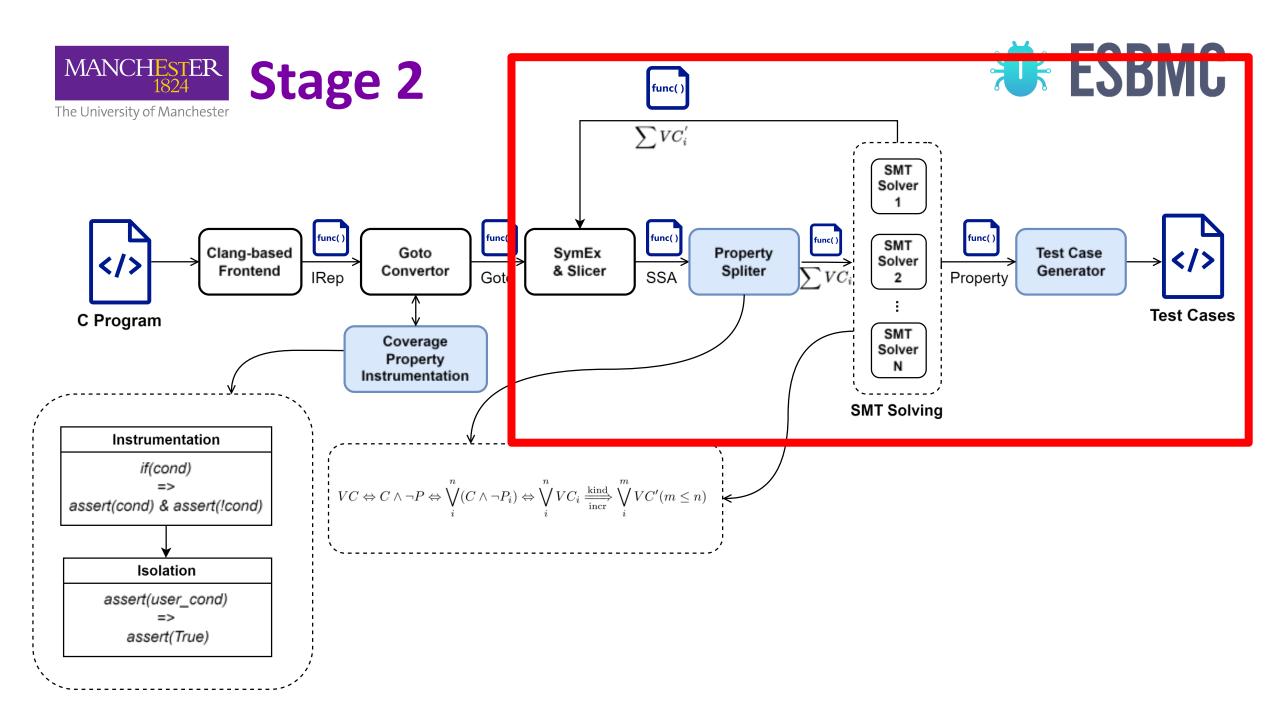
Internal safety checks within ESBMC are disabled ۲

> /* assert(y!=0) */ Double z = x / y













func() $\sum VC'_i$ SMT Solver ٠ func() func() func() SMT Property SymEx Solver & Slicer Spliter VC_i SSA Goto Property 2 SMT Solver SMT Solving

- The CFG get symbolic executed within defined bounds (e.g., through loop unrolling) and is eventually encoded as a **verification condition (VC)**.
- In ESBMC, the VC is an SMT formula incorporating:
 - Constraints (execution conditions, C)
 - Properties (expected behaviours, P)

 $VC \Leftrightarrow C \wedge \neg P$

• To let ESBMC verify multiple properties, we split the property *P* into a set of unit properties *P_i*

$$C \wedge \neg P \Leftrightarrow \bigvee_{i}^{n} (C \wedge \neg P_{i}) \Leftrightarrow \bigvee_{i}^{n} VC_{i}$$





Issue:

- Normally, BMC verifies system behavior only up to a fixed bound *k*, once this threshold is reached, it terminates. The verification result becomes unknown.
- As a consequence, some branch paths may be missed.

How about we set a relatively large *k* (e.g. *100*)?

- No guarantee of soundness: Larger *k* increases depth but doesn't ensure that all paths are covered or that unreachability is proved.
- Inefficiency: Large bounds might lead to state-space explosion





Aid: Use incremental reasoning (e.g. k-induction) to automatically extend verification beyond bound k

- > It checks that a property holds up to a bound (k)
- > It proves that if it holds at (k), it also holds at (k + 1)

In ESBMC, this can be summaries as three steps:

- **Base Case**: Check if the program is correct up to (*k*) steps (normal BMC).
- Forward Reasoning: prove that all loops in the program were fully unrolled.

B(k)

 \wedge

I(k)

• Inductive Step: If it is good up to (k), prove it's still good for (k + 1).

 $\neg \quad B(k) \quad o \quad ext{program contains bug}$

 \rightarrow program is correct

 $B(k) \ \wedge \ F(k) \ o \ ext{program is correct}$

https://ssvlab.github.io/esbmc/documentation.html#k-induction

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```
Bool return_value VERIFIER_nondet_bool_1;
return value VERIFIER nondet bool 1 = VERIFIER nondet bool();
assert(!return_value____VERIFIER_nondet_bool_1); // !return_value$___VERIFIER_nondet_bool$1
assert(!(!return value VERIFIER nondet bool 1)); // !(!return value$ VERIFIER nondet bool$1)
if (!return value VERIFIER nondet bool 1)
   goto ESBMC goto label 4;
```

```
assert(!(!( Bool)x)); // !(!( Bool)x)
assert(!(!(!( Bool)x))); // !(!(!( Bool)x))
if (!(!( Bool)x))
   goto ESBMC goto label 2;
```

```
assert(!(x == 1)); // !(x == 1)
assert(!(!(x == 1))); // !(!(x == 1))
if (!(x == 1))
   goto ESBMC goto label 3;
```





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<pre>!(!return_value\$VERIFIER_nondet_bool\$1)</pre>
<pre>!return_value\$VERIFIER_nondet_bool\$1</pre>
!(!(!(_Bool)x))
!(!(_Bool)x)
!(!(x == 1))
!(x == 1)

Termination:

a) all remaining coverage properties are proven during forward reasoning

b) all properties are reduced to tautologies and removed through slicing, leaving no properties for further verification **Re-verification**: if any property P_i remains unknown, a verification re-run is initiated



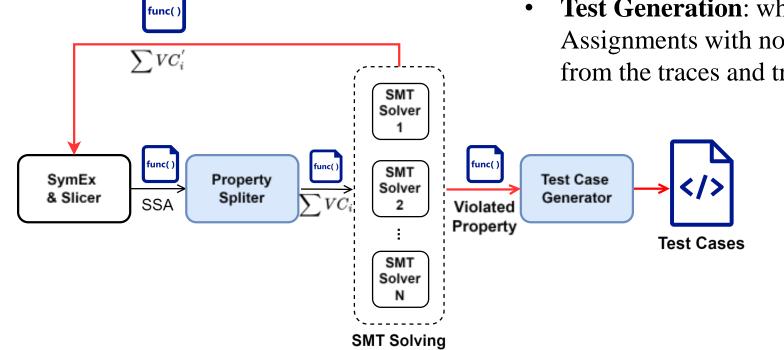
$$\bigvee_{i}^{n} VC_{i} \xrightarrow{\text{kind}}_{\text{incr}} \bigvee_{i}^{m} VC'(m \le n)$$





Solution found by the inductive step (k = 4)
[Coverage]
Branches : 6
Reached : 6
 !(!(!(_Bool)x)) file exp.c line 11 column 5 function main
 !(x == 1) file exp.c line 16 column 10 function main
 !return_value\$___VERIFIER_nondet_bool\$1 file exp.c line 9 column 3 function main
 !(!return_value\$___VERIFIER_nondet_bool\$1) file exp.c line 9 column 3 function main
 !(!(x == 1)) file exp.c line 16 column 10 function main
 !(!(_Bool)x) file exp.c line 11 column 5 function main
 !(!(_Bool)x) file exp.c line 11 column 5 function main
 Branch Coverage: 100%

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• **Test Generation**: whenever a property P_i violation is reported. Assignments with nondeterministic initial values are extracted from the traces and transformed into corresponding test suites.

testcase-0.xml
testcase-1.xml
testcase-2.xml
testcase-3.xml
testcase-4.xml
testcase-5.xml





ESBMC ranked around 7th–8th among all participants in Cover-Branches category at Test-Comp 2025.

Participants	Plots	<u>cetfuzz</u>	<u>CoVeriTes</u>	ESBMC- incr	ESBMC- kind	FDSE	<u>Fizzer</u>	FuSeBMC	FuSeBMC- Al	<u>HybridTiger</u>	KLEE	KLEEF	<u>Owi</u>	PRTest	Rizzer	<u>Sikraken</u>	<u>Symbiotic</u>	<u>TracerX</u>	TracerX- WP	<u>UTestGen</u>	WASP-
Cover-Branches 10011 valid tasks	I Grandle Bla	<u>2524</u>	<u>4959</u>	<u>4380</u>	<u>4323</u>	<u>5468</u>	<u>5429</u>	<u>5656</u>	<u>4077</u>	<u>3866</u>	<u>3065</u>	<u>5734</u>	<u>2462</u>	<u>3191</u>	_	<u>2469</u>	<u>4207</u>	<u>3327</u>	<u>3275</u>	<u>4393</u>	<u>2740</u>
CPU time	·	1800000 s	6000000 s	1500000 s	980000 s	7900000 s	7300000 s	8900000 s	3600000 s	4800000 s	2700000 s	5400000 s	6200000 s	8200000 s		7300000 s	2800000 s	2000000 s	2300000 s	5900000 s	3600000 s

Here, ESBMC-incr performs slightly better, mostly due to different unwinding configuration.





ESBMC is open-source under the Apache License 2.0, and its C++ source code is publicly available on GitHub: <u>https://github.com/esbmc/esbmc/</u>

The official website is available at: <u>https://esbmc.org</u>

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.github/workflows	Fix windows build (#2403)	last week	€ esbmc.org/									
docs	[Documentation] Added examples fro	3 months ago	kotlin python c cpp smt-solver									
regression	[python-frontend] Handling list slicin	9 hours ago	k-induction automated-testing incremental-learning bmc									
scripts	Update stats-30s.txt	last week	solidity-contracts cheri									
src	[python-frontend] Handling list slicin	9 hours ago	automated-verification cp-solver									





Thank you!