Neural Network Verification is a Programming Language Challenge

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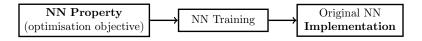
# The textbook ML workflow



## Three ingredients

- Property. Expressed as a loss function
- > Training. Empirical risk minimisation over finite data
- Implementation. ML frameworks, HW accelerators

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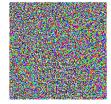
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57.7% confidence

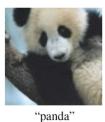
 $+.007 \times$ 





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"gibbon" 99.3 % confidence



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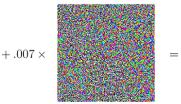
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## Vision

"Please make the neural network safe"







"panda" 57.7% confidence

"gibbon" 99.3 % confidence

#### Local robustness property

- For each image x in the training set
- ▶ Define a small local set of perturbations  $||x x'|| \le \varepsilon$
- Ensure output class is arg max{f(x)} = arg max{f(x')}



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=

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### Programming languages interpretation

Refinement type for f with predicate

$$\blacktriangleright \quad \forall x, ||x - x'|| \le \varepsilon \implies ||f(x) - f(x')|| \le \delta$$

# Other "desirable" properties

### A high-level taxonomy

Geometric properties. Defined along the data manifold. Examples: local robustness and equivalence, fairness.



"panda" 57.7% confidence





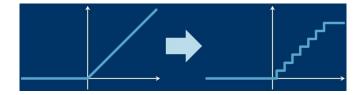
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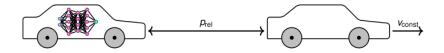
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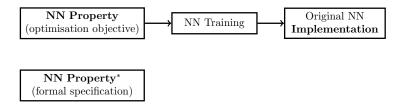
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- Hyper-properties. Must hold for any inputs.
  Examples: global robustness and equivalence, monotonicity.
- Domain specific. Based on semantic of ML task. Examples: arbitrary pre- and post-conditions.





Vision (v2)

"Please check whether the neural network is safe"

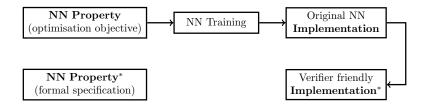


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### Solution

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# Vision (v2)



"Please check whether the neural network is safe"

#### Solution

- Property\*. Define what "safe" means.
- **Implementation\*.** Abstract away SW/HW details.
- **Verification.** Run a state-of-the-art verifier.



# Existing verification pipeline (1)



## Objective

- Given a neural network  $f : \mathbb{R}^n \to \mathbb{R}^m$
- Prove that property  $\Xi(f)$  holds

<sup>&</sup>lt;sup>1</sup>International competition: VNN-COMP (since 2020)

# Existing verification pipeline (1)



## Objective

- Given a neural network  $f : \mathbb{R}^n \to \mathbb{R}^m$
- Prove that property  $\Xi(f)$  holds

## State of the art

- f is expressed as ONNX or other exchange format
- ► Tools: Marabou,  $\alpha\beta$ -CROWN, PyRAT, NNV, ERAN...
- accept  $\Xi(f)$  as **linear** pre- and post-conditions on f
- usually expressed in VNN-LIB format<sup>1</sup>

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# Challenges

- Cannot write properties on two networks  $f_1, f_2$
- Properties must list all I/O entries explicitly
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- Cannot automatically bind to dataset entries
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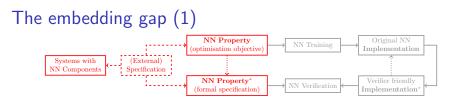
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## Existing solutions

- CAISAR: WhyML-like specification language
- Vehicle: dependent-typed functional language



## Objective

- We want to prove a **system** property  $\Psi(s(\cdot))$
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## Strategy

• We decompose it as  $\Xi(f) \implies \Phi(u \circ f \circ e) \implies \Psi(s(u \circ f \circ e))$ 

### where Ξ(f) is a property of the neural network alone Φ(u ∘ f ∘ e) maps it back to the problem space

# The embedding gap (2)

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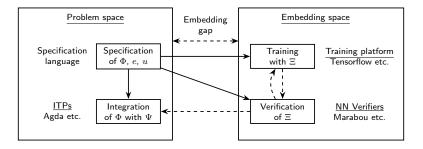
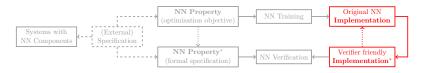
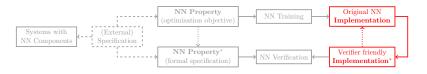


Figure: Outline of Vehicle compiler backends.



### Mismatch in numerical types

- Original implementation  $\rightarrow$  **finite** precision
- $\blacktriangleright \text{ Verified implementation} \rightarrow \textbf{real-valued} \text{ arithmetic}$

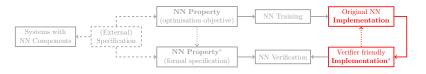


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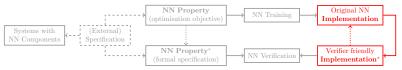
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#### Quantised neural networks

Require reasoning about integer arithmetic



### Many sources of non determinism

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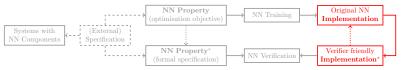
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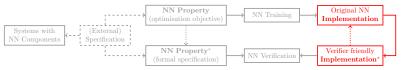
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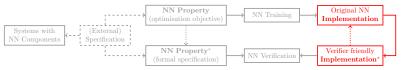
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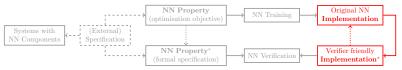
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#### Access to software implementation is required!

# Proof production and checking



### ONNX operator support

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- E.g. abstract interpretation procedure for new activations
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## Safety proof certificates

- Safety proofs are more difficult to find than counterexamples
- Format: usually some form of branch-and-bound tree
- Challenge 1: formalise under Farkas' lemma
- Challenge 2: write verified proof checkers

# Property-guided training



#### ML-like approaches

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- ▶ Differentiable logic: property  $\rightarrow$  loss function

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#### Certified approaches

- Interval Bound Propagation (IBP) for robustness training
- Lipschitz-bounded, monotonic & convex neural networks
- Certified training on a limited number of other properties

# Verification of cyber-physical systems



## CPS

- Feedback controller for plant model (ODEs)
- Requires hybrid description (discrete + continuous)
- International competition: AINNCS category @ ARCH-COMP
- Tools: CORA, JuliaReach, NNV, OVERT, POLAR...
- Challenges: property specification, scalability, software...

# Neural network verification today

Existing Solutions	High-level		Low-level		Quantised		Software		Future	
PL Challenges	Vehicle	CAISAR	lphaeta-CROWN	Marabou	QEBVerif	Aster	CBMC	ESBMC	Unified Language	Formal Interfaces
Rigorous Semantics	$\checkmark$	$\checkmark$					$\checkmark$	$\checkmark$	✓	$\checkmark$
Embedding Gap Implementation Gap Proof Certificates Supports Training	$\checkmark$	$\checkmark$		$\checkmark$	$\checkmark$	$\checkmark$	√ √*	√ √*		√ √ √

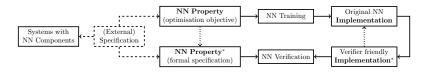
#### Five desirable PL features

- Current tools  $\rightarrow$  **partial** support
- ▶ We need a more principled solution





# The future roadmap (1)



### Option 1: a unified language

- Rigorous semantics. Use dependent types. All components are implemented in a signle language.
- **Embedding gap.** Becomes a type conversion problem.
- Implementation gap. Enforce explicit types. No external libraries, unless formally verified.
- **Proof certificates.** The type checker is the proof checker.
- Training support. Requires re-implementation or code synthesis in the new language.

# The future roadmap (2)

## Option 2: formal interfaces

- More **flexible**: can accommodate existing frameworks.
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- Keep a maximally-expressive specification language.
- Design a compiler that can use existing tools.
- Each tool solves a specific verification/synthesis sub-goal.
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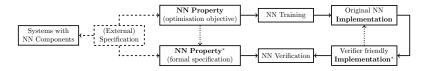
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## Our inspiration

 Behavioural interface specification languages (BISL) such as JML, Why3, SPARK

# A diagrammatic summary



Existing Solutions	kisting Solutions High-level		Low-level		Qua	Quantised		Software		Future	
PL Challenges	Vehicle	CAISAR	lphaeta-CROWN	Marabou	QEBVerif	Aster	CBMC	ESBMC	Unified Language	Formal Interfaces	
Rigorous Semantics Embedding Gap	√ √	<b>√</b>					<b>√</b>	~	√ ✓	√ √	
Implementation Gap Proof Certificates Supports Training	$\checkmark$	$\checkmark$		$\checkmark$	$\checkmark$	$\checkmark$	√ √*	√ √*		√ √	

Any questions?